ADT: Design & Implementation

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Outline

- Concept: ADT
- Demonstration of ADT's design and implementation
 - List as an ADT (our focus)
 - Stack & Queue as ADTs (exercises)

Abstract Data Type (ADT)

- ADT = data items + operations on the data
- Design of ADT

- Determine data members & operations

- Implementation of ADT
 - Storage/data structures to store the data

- Algorithms for the operations, i.e., how to do it

Example: List

- List as an ADT
- 3 different implementations
 - Static array-based
 - Dynamic array-based
 - Linked-list based

Properties of Lists

- Can have a single element
- Can have <u>no</u> elements \rightarrow empty!
- There can be lists of lists
- We will look at the list as an abstract data type
 - Homogeneous
 - Finite length
 - Sequential elements

Basic Operations

- Construct an empty list
- Determine whether or not empty
- Insert an element into the list
- Delete an element from the list
- Traverse (iterate through) the list to
 - Modify
 - Output
 - Search for a specific value
 - Copy or save
 - Rearrange

Designing a List Class

- Should contain at least the following function members
 - Constructor
 - empty()
 - insert()
 - delete()
 - display()
- Implementation involves
 - Choosing data structures to store data members
 - Choosing algorithms to implement function members

Impl. 1: Static Array

- Use a static array to hold list elements
- Easy to manipulate arrays for implementing operations
- Natural fit for mapping list elements to array elements



Design

- What data members?
- What operations?

Design

```
const int CAPACITY = 1000; //for what purpose?
typedef int ElementType; //for what purpose?
```

```
class List {
    private:
        int size; //# of elements
        ElementType array[CAPACITY];
    public:
```

};

...

Implementing Operations

- Constructor
 - Static array allocated at compile time. No need to allocate explicitly!
- Empty
 - Check if size == 0
- Traverse
 - Use a loop from 0th element to size 1
- Insert
 - Shift elements to right of insertion point
- Delete
 - Shift elements back



List Class with Static Array - Problems

- Stuck with "one size fits all"
 - Could be wasting space
 - Could run out of space
- Better to have instantiation of specific list: specify what the capacity should be
- Thus we consider creating a List class with dynamically-allocated array

Impl. 2: Using Dynamic Arrays

• Allow List objects to specify varied sizes

Dynamic-Allocation for List Class

Now possible to specify different sized lists
 cin >> maxListSize;
 List aList1(maxListSize);
 List aList2(500);



Dynamic-Allocation for List Class

- Changes required in data members
 - Eliminate const declaration for CAPACITY
 - Add data member to store capacity specified by client program
 - Change array data member to a pointer
 - Constructor requires considerable change
- Little or no changes required for
 - empty()
 - display()
 - erase()
 - insert()

Design

```
typedef int ElementType;
```

class List {

private:

int mySize; //# of elements
int myCapacity;
ElementType *myArray;

public:

```
List (int cap);
```

};

What needs to be changed?

- Constructor
- Addition of other functions to deal with dynamically allocated memory
 - Destructor
 - Copy constructor (won't be discussed here)
 - Assignment operator (won't be discussed here)

Constructor

- Two tasks
 - Data member initialization
 - Dynamically allocate memory

//constructor

}

List::List(int cap) : mySize(0), myCapacity(cap) { myArray = new int[cap]; //allocate memory assert(myArray != 0); //need #include <cassert>

Destructor Needed!

- When class object goes out of scope, the pointer to the dynamically allocated memory is reclaimed automatically
- The dynamically allocated memory is not



• The destructor reclaims dynamically allocated memory (the default destructor fails here!)

Destructor

- The default destructor needs to be overrided!
- De-allocate memory you allocated previously; otherwise causes memory leak problem!

```
List::~List() {
delete [] myArray;
```

Problems

- Array capacity cannot be changed during running time (execution)!
- Insertion & deletion are not efficient!

Involve element shifting

Impl. 3: Using Linked-Lists

 To address the two problems faced by arraybased solutions

Review: Linked List

• Linked list nodes contain

fin

- Data part - stores an element of the list

data

REXI

Next part – stores link/pointer to next element
 (when no next element, null value)

Basic Operations

- Traversal
- Insertion
- Deletion

Traversal

• Initialize a variable ptr to point to first node



Process data where ptr points

Traversal (cont.)

- set ptr = ptr->next, process ptr->data
 first 9
 first 9
 ptr 17
 22
 26
 34
 ptr 17
 22
 26
 34
- Continue until ptr == null

Insertion



- Insertion
 - To insert 20 after 17
 - Need address of item before point of insertion
 - predptr points to the node containing 17
 - Create a new node pointed to by newptr and store 20 in it
 - Set the next pointer of this new node equal to the next pointer in its predecessor, thus making it point to its successor.
 - Reset the next pointer of its predecessor to point to this new node

Insertion

- Note: insertion also works at <u>end</u> of list
 pointer member of new node set to null
- Insertion at the <u>beginning</u> of the list
 <u>predptr</u> must be set to <u>first</u>
 - pointer member of newptr set to that value
 - first set to value of newptr

Note: In all cases, **no shifting of list** elements is required !



• Delete node containing 22 from list.

To free space

- Suppose **ptr** points to the node to be deleted
- predptr points to its predecessor (the 17)
- Do a bypass operation:
 - Set the next pointer in the predecessor to point to the successor of the node to be deleted
 - Deallocate the node being deleted.

1 Lesson in Pointers

Before operating on an pointer, first check if it is null!

if (ptr) {
 //do sth. meaningful
}

Linked Lists - Advantages

- Access any item as long as external link to first item maintained
- Insert new item without shifting
- Delete existing item without shifting
- Can expand/contract as necessary

Linked Lists - Disadvantages

- Overhead of links:
 - used only internally, pure overhead
- If dynamic, must provide
 - destructor
 - copy constructor (but not here!)
- No longer have direct access to each element of the list
 - Many sorting algorithms need direct access
 - Binary search needs direct access
- Access of nth item now less efficient
 - must go through first element, and then second, and then third, etc.

Linked Lists - Disadvantages

- List-processing algorithms that require fast access to each element cannot be done as efficiently with linked lists.
- Consider adding an element at the end of the list

Array	Linked List
a[size++] = value;	Get a new node;
	set data part = value
	next part = <i>null_value</i>
	If list is empty
	Set first to point to new node.
	Else
	Traverse list to find last node
This is the inefficient part	Set next part of last node to
	point to new node.

Using C++ Pointers and Classes

• To Implement Nodes

```
class Node
{
    public:
    DataType data;
    Node * next;
};
```

- Note: The definition of a Node is <u>recursive</u>
 (or self-referential)
- It uses the name Node in its definition
- The next member is defined as a pointer to a Node

Working with Nodes

• Declaring pointers

Node* ptr; Or typedef Node* NodePointer; NodePointer ptr;

- Allocate and deallocate
 ptr = new Node; delete ptr;
- Access the data and next part of node (*ptr).data and (*ptr).next or

ptr->data and ptr->next

Working with Nodes

 Note data members are public



 This class declaration will be placed inside another class declaration for List (private section), p296

Implementing List with Linked-Lists



typedef Node * NodePointer;

Design & Impl.

- Add data members
- Add member functions

Class List

```
#ifndef _LIST_H
#define _LIST_H
typedef int ElementType;
class List {
  private:
     class Node {
        public:
           ElementType data;
           Node* next;
     };
     typedef Node* NodePointer;
  private:
    NodePointer first;
  public:
     List();
     ~List();
     void insert(ElementType x);
     void delete(ElementType x);
};
#endif
```

Implementing

- Constructor
- Destructor
- Operation: insert()
- Operation: delete()

Exercises

• Implement Stack and Queue classes with different implementations.

Stack

- Design and Implement a Stack class
- 3 options
 - Static array
 - Dynamic array
 - Linked list

Stack.h

```
typedef int DataType;
class Stack {
   public:
      Stack();
      Stack(const Stack& org);
      void push(const DataType& v);
      void pop();
      DataType top() const;
      ~Stack();
   private:
      class Node {
        public:
           DataType data;
           Node* next;
           Node(DataType v, Node* p) : data(v), next(0) { }
      };
      typedef Node* NodePtr;
      NodePtr myTop;
};
```

Queue

- Design and Implement a Queue class
- 3 options
 - Static array
 - Dynamic array
 - Linked list

Queue.h

```
typedef int DataType;
class Queue {
   public:
     //constructor
    //... member functions
   private:
     class Node {
        public:
          DataType data;
          Node* next;
          Node(DataType v, Node* p) : data(v), next(0) { }
      };
      typedef Node* NodePtr;
     NodePtr myFront, myback;
```

};