

# Monic representations for higher-rank graphs

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joint with C. Farsi, P. Jorgensen, S. Kang, J. Packer

West Coast Operator Algebras Seminar  
Seattle University, 7 October 2018

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- (2015) For higher-rank graphs  $\Lambda$ , Farsi–G–Kang–Packer introduced  $\Lambda$ -semibranching function systems & the associated reps.

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- (2017/18) Farsi–G–Jorgensen–Kang–Packer characterize all monic representations of  $C^*(\Lambda)$  in terms of  $\Lambda$ -SBFS. In particular, since Cuntz–Krieger algebras are higher-rank graph algebras, this answers the question posed by Bezuglyi–Jorgensen.

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- “Measure-theoretic” representations:  $\Lambda$ -semibranching function systems
- $\Lambda$ -projective systems
- Monic representations; main Theorem.

# Higher-rank graphs

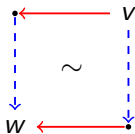
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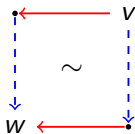
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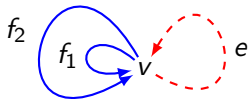
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Introduced by Kumjian & Pask in 2000 to give examples of combinatorial, computable  $C^*$ -algebras, more general than  $C^*(E)$ .

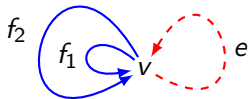
Paths in  $E \rightsquigarrow k$ -dimensional rectangles in  $\Lambda$ .

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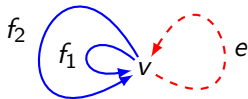
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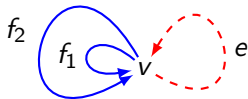
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Note that  $\Lambda^0$  is the vertices of  $\Lambda$ .

# Higher-rank graph $C^*$ -algebras

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$$(CK4) \quad \text{For any } v \in \text{Obj } \Lambda \text{ and any } n \in \mathbb{N}^k,$$

$$p_v = \sum_{\lambda \in v\Lambda^n} s_\lambda s_\lambda^*.$$

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Moreover,  $s_e s_{f_1} = s_{f_2} s_e$ , so the action of  $\mathbb{Z}$  on  $\mathcal{O}_2 = C^*(s_{f_1}, s_{f_2})$  is given by

$$\alpha(s_{f_i}) = s_{f_{i+1}}.$$

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$$\tau_\lambda : D_{s(\lambda)} \rightarrow R_\lambda, \quad \tau^n : R_\lambda \rightarrow D_{s(\lambda)}$$

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- Various rules (analogous to (CK1)-(CK4) ) governing how these maps and sets interact.
- Positive a.e. Radon–Nikodym derivatives

$$\Phi_\lambda := \frac{d(\mu \circ \tau_\lambda)}{d\mu}.$$

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For our nontrivial example  $\Lambda$ , we can define two  $\Lambda$ -SBFS's on Lebesgue measure spaces: one on  $X = [0, 1]$  and one on  $Y = [0, 1]^2$ .

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This latter example is the only one I know of that has a non-constant Radon-Nikodym derivative.

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$\Lambda$ -projective systems allow for a more general scaling function.

# $\Lambda$ -projective representations

Assume we have a  $\Lambda$ -SBFS on  $(X, \mu)$ .

Theorem (Farsi-G-Jorgensen-Kang-Packer 2017)

Suppose a family  $\{f_\lambda : \lambda \in \Lambda\} \subseteq L^2(X, \mu)$  of functions, with  $f_\lambda$  supported on  $R_\lambda$ , satisfies

$$f_\lambda \cdot (f_\nu \circ \tau^{d(\lambda)}) = f_{\lambda\nu} \quad \text{and} \quad |f_\lambda|^2 = \left(\Phi_\lambda \circ \tau^{d(\lambda)}\right)^{-1} \Big|_{R_\lambda}.$$

Then the operators  $\{T_\lambda\}_{\lambda \in \Lambda}$  give a representation of  $C^*(\Lambda)$  on  $L^2(X, \mu)$ :

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This gives the  $\Lambda$ -SBFS representation:  $T_\lambda = \pi(s_\lambda)$ .

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Monic systems always give a representation of the Cuntz(–Krieger) algebra.

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Monic systems always give a representation of the Cuntz(–Krieger) algebra.

## Corollary

For Cuntz and Cuntz–Krieger algebras,  $\Lambda$ -projective systems and monic systems are the same thing.

# Infinite paths in higher-rank graphs

In order to define monic representations we need to discuss  $\Lambda^\infty$ , the space of infinite paths in  $\Lambda$ .

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The sets  $Z(\lambda)$ , for  $\lambda \in \Lambda$ , form a compact open basis for the topology on  $\Lambda^\infty$ :

$$Z(\lambda) = \{x \in \Lambda^\infty : x = \lambda y \text{ for some } y \in \Lambda^\infty\}.$$

# Infinite paths and $\Lambda$ -SBFS

When  $\Lambda$  is finite and strongly connected, we have a canonical measure  $M$  on  $\Lambda^\infty$ , due to an Huef, Laca, Raeburn, and Sims.

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## Theorem (Farsi-G-Kang-Packer, 2015)

*Any finite, strongly connected  $k$ -graph admits a  $\Lambda$ -SBFS on the measure space  $(\Lambda^\infty, M)$ :*

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$$\phi_\lambda = \frac{M(Z(\lambda))}{M(Z(s(\lambda)))} \chi_{Z(s(\lambda))}.$$

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## Theorem (Farsi-G-Jorgensen-Kang-Packer, 2017)

*Let  $\Lambda$  be a finite  $k$ -graph with no sources. A representation  $\phi$  of  $C^*(\Lambda)$  is monic iff  $\phi$  is unitarily equivalent to a  $\Lambda$ -projective representation on  $(\Lambda^\infty, \mu)$  where  $\mu$  is a Borel measure on  $\Lambda^\infty$ .*

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## Corollary

*Any  $\Lambda$ -SBFS on  $\Lambda^\infty$  gives a monic representation. Perhaps more than one.*



## Theorem (Farsi-G-Jorgensen-Kang-Packer)

*Let  $\Lambda$  be a finite  $k$ -graph with no sources. Suppose we have a  $\Lambda$ -SBFS on  $(X, \mu)$ . The associated  $\Lambda$ -semibranching representation  $\pi : C^*(\Lambda) \rightarrow B(L^2(X, \mu))$  is monic iff*

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It also answers the question of Bezuglyi and Jorgensen:

For Cuntz–Krieger algebras, the monic systems which give rise to monic representations are precisely the ones associated to SBFS where the range sets  $R_i$  generate the  $\sigma$ -algebra.

# The end

Thanks for your attention!